

Monads

October 28, 2019

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A burrito

Monads

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Burritos

- Monads are like burritos

Burritos

- Monads are like burritos
- Monads are not like burritos

Sequencing Actions

1. Get a line
2. Get a line
3. “Return” the lines concatenated together

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 - `myAction = do`
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 `return $ a ++ b`

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  b <- getLine
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```
  return $ a ++ b
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```
    = (++) <$> getLine <*> getLine
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Sequencing Actions

1. Get a line
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3. **Print** the lines concatenated together

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 - `myAction = do`
 - `a <- getLine`
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 - `print $ a ++ b`
 - How to write this in applicative style?

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- Actions

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- What to do with the results

Sequencing Actions

1. Get a line
2. Get a line
3. Print the lines concatenated together

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```
myAction' = (\x y -> print $ x ++ y)
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```

Sequencing Actions

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 `a <- getLine`
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 `print $ a ++ b`

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myAction' = (\x y -> print $ x ++ y)
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- Why doesn't this work?

Sequencing Actions

- `(\x y -> print $ x ++ y) <$> getLine <*> getLine`
 - Get a line `a`, apply `(\x y -> print $ x ++ y)` to `a` (to get `(\y -> print $ a ++ y)`), and wrap it up in an IO box

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 - We never actually ran `print $ a ++ b`!
- `myAction :: IO ()`
- `myAction' :: IO (IO ())`

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 - Take `(\y -> print $ a ++ y)` out of the box, get another line `b`, apply `(\y -> print $ a ++ y)` to `b` (to get `print $ a ++ b`), and **wrap it up in another IO box**
 - We never actually ran `print $ a ++ b`!
- `myAction :: IO ()`
- `myAction' :: IO (IO ())`
 - To run `print $ a ++ b`, we need to take it out of the box

Monads

- Wikipedia: Throughout this article C denotes a category. A monad on C consists of an endofunctor $T: C \rightarrow C$ together with two natural transformations:
 $\eta: 1_C \rightarrow T$ (where 1_C denotes the identity functor on C) and
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 - Remember categories:

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 - Remember categories:
 - category = objects + morphisms
 - objects = types
 - morphisms = functions

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 - Our only category is Hask, so all functors are endofunctors

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 - natural transformation = morphism of functors
 - Let us call η unit (or return), and μ join
 - If Haskell syntax allowed it, we could say
`return :: Identity -> T` and
`join :: T2 -> T`

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- Throughout this article C denotes a category. A monad on C consists of an endofunctor T together with two natural transformations:
`return :: a -> T a` and
`join :: T (T a) -> T a.`

Sequencing Actions

- `myAction' :: IO (IO ())`

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Sequencing Actions

- `myAction' :: IO (IO ())`
- `join myAction' :: IO ()`
- `Prelude Control.Monad> join myAction'`
`the_`
`dog`
`"the_dog"`

Monads

- class Monad m where

```
    return :: a -> m a
```

```
    (>>=) :: m a -> (a -> m b) -> m b
```

```
    (>>) :: m a -> m b -> m b
```

```
x >> y = x >>= \_ -> y
```

```
fail :: String -> m a
```

```
fail msg = error msg
```

Monads

- `class (Applicative m) => Monad m` where
`return :: a -> m a`

`(>>=) :: m a -> (a -> m b) -> m b`

`(>>) :: m a -> m b -> m b`

`x >> y = x >>= _ -> y`

`fail :: String -> m a`

`fail msg = error msg`

- Since GHC v7.10, `Applicative` is a superclass of `Monad`

Monads

- `class (Applicative m) => Monad m` where
`return :: a -> m a`

`(>>=) :: m a -> (a -> m b) -> m b`

`(>>) :: m a -> m b -> m b`

`x >> y = x >>= _ -> y`

`fail :: String -> m a`

`fail msg = error msg`

- What happened to `join`? What are `(>>=)`, `(>>)`, and `fail` doing here?

Monads

- $(\gg=) :: m\ a \rightarrow (a \rightarrow m\ b) \rightarrow m\ b$

Monads

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- $(= <<) = \text{flip } (\gg=)$

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$(= <<) :: (a \rightarrow m\ b) \rightarrow m\ a \rightarrow m\ b$

- $(\langle * \rangle) :: f\ (a \rightarrow b) \rightarrow f\ a \rightarrow f\ b$

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$(=<<) :: (a \rightarrow m\ b) \rightarrow m\ a \rightarrow m\ b$

- $(<*>) :: f\ (a \rightarrow b) \rightarrow f\ a \rightarrow f\ b$
- $(<\$>) :: (a \rightarrow b) \rightarrow f\ a \rightarrow f\ b$

- $(=<<)$ (and $(>>=)$) are maps for **monadic functions**

Monads

- $(\gg=) :: m\ a \rightarrow (a \rightarrow m\ b) \rightarrow m\ b$
- $(= <<) = \text{flip } (\gg=)$

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- $(\langle * \rangle) :: f\ (a \rightarrow b) \rightarrow f\ a \rightarrow f\ b$
- $(\langle \$ \rangle) :: (a \rightarrow b) \rightarrow f\ a \rightarrow f\ b$

- $(= <<)$ (and $(\gg=)$) are maps for **monadic functions**
 - Functions that create their own boxes

Monads

- $(\gg=) :: m\ a \rightarrow (a \rightarrow m\ b) \rightarrow m\ b$
- $(= \ll) = \text{flip } (\gg=)$

$(= \ll) :: (a \rightarrow m\ b) \rightarrow m\ a \rightarrow m\ b$

- $(\ll*) :: f\ (a \rightarrow b) \rightarrow f\ a \rightarrow f\ b$
- $(\ll\$) :: (a \rightarrow b) \rightarrow f\ a \rightarrow f\ b$

- $(= \ll)$ (and $(\gg=)$) are maps for **monadic functions**
 - Functions that create their own **context**

Monads

- $g \gg= f = \text{join } (\text{fmap } f \ g) :: m \ a \ \rightarrow \ (a \ \rightarrow \ m \ b) \ \rightarrow \ m \ b$

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 - $\text{fmap } f$ lifts it to type $m \ a \ \rightarrow \ m \ (m \ b)$

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 - $g :: m \ a$ is a value of type a in a box
 - $\text{fmap } f \ g :: m \ (m \ b)$ outputs a value of type b in two nested boxes

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 - $\text{fmap } f \ g :: m \ (m \ b)$ outputs a value of type b in two nested boxes
 - $\text{join } (\text{fmap } f \ g)$ extracts a monadic value of type $m \ b$ from the outermost box

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 - $g \gg= f$ extracts a value of type a from g and feeds it to f to get a monadic value of type $m \ b$
- $\text{join } x = x \gg= \text{id}$

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- `class (Applicative m) => Monad m` where
`return :: a -> m a`

`(>>=) :: m a -> (a -> m b) -> m b`

`(>>) :: m a -> m b -> m b`

`x >> y = x >>= _ -> y`

`fail :: String -> m a`

`fail msg = error msg`

- Shorthand for when we don't need to bind the value inside `x` to evaluate `y`

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- Error handler for pattern matching in `do` expressions